

# riscure

## **Efficient practical key recovery for side channel attacks**

Ilya Kizhvatov

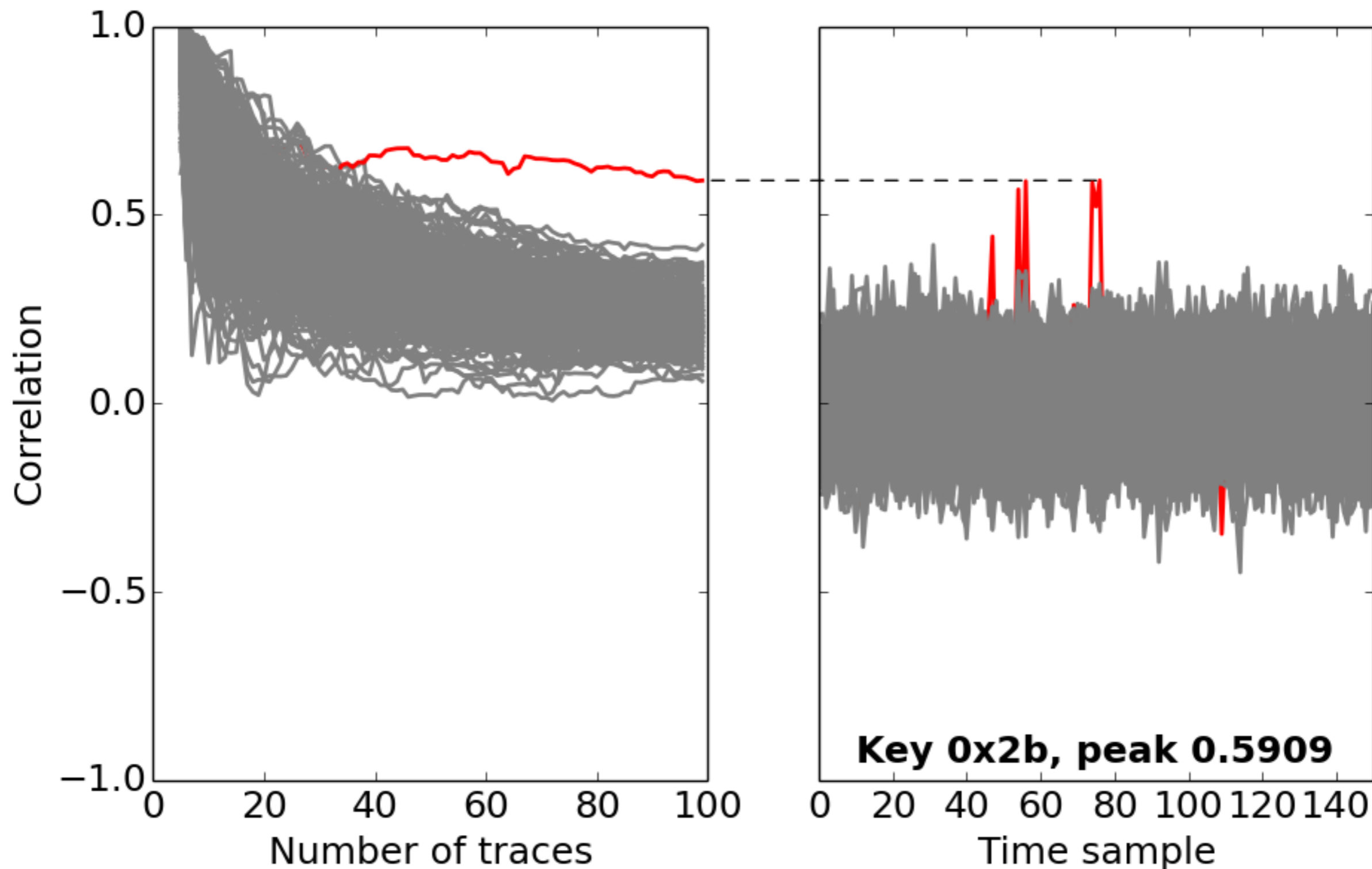
Joint work with Andrey Bogdanov and Kamran Manzoor (DTU), and Marc Witteman (Riscure)

MCrypt, Les Deux Alpes, 15 August 2014

# DPA recap

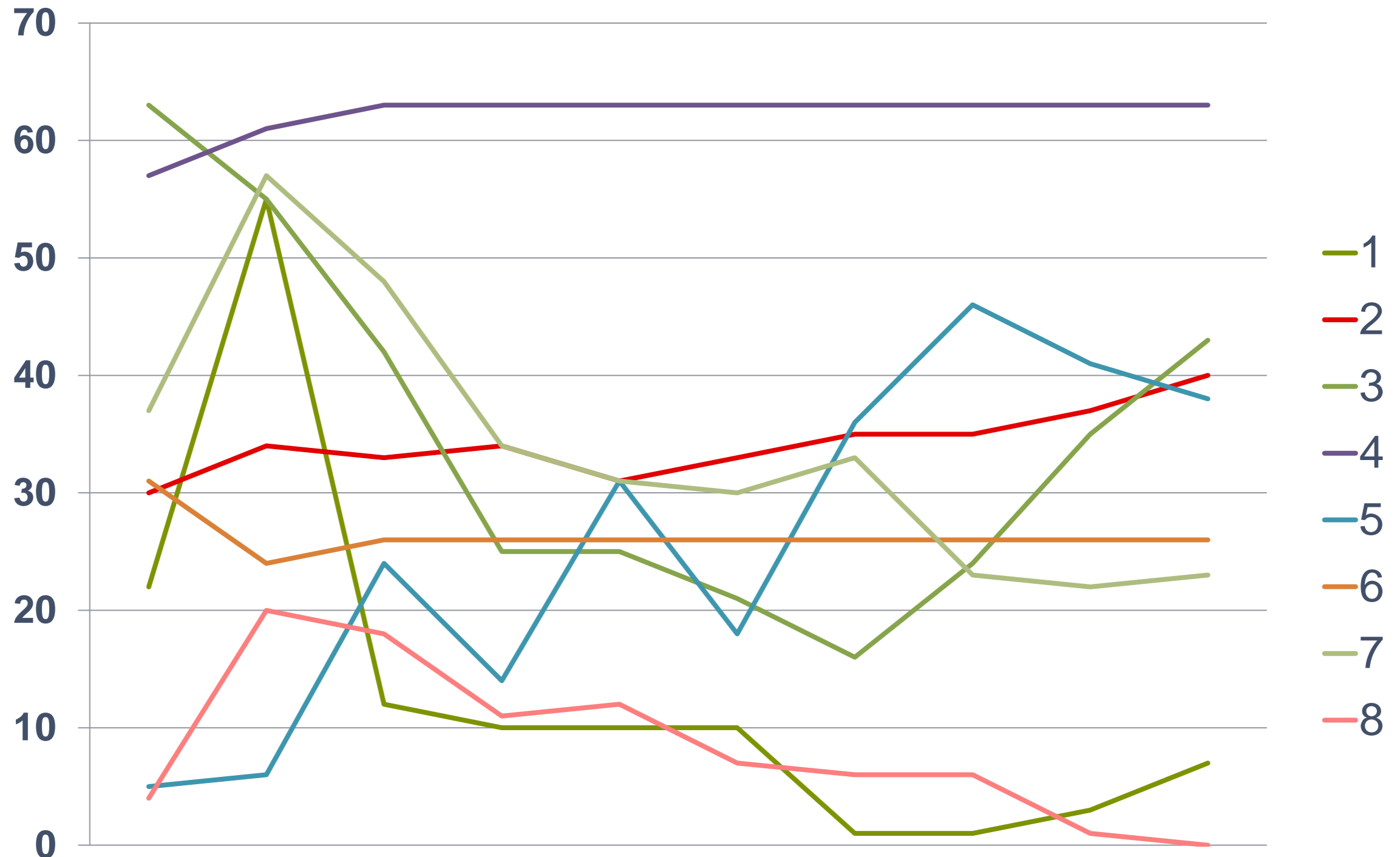
Distinguishing key chunks

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# DPA recap

Evolution of the rank of correct key chunk candidates



# Problem statement

## Conquer phase:

- The sub-keys are combined to generate the full-key.
- *Ideally* full-key = the most probable sub-key candidate from each list

$k_0$	0.0065*	$k_0$	0.0071	.....	$k_0$	0.0070
$k_1$	0.0063	$k_1$	0.0068		$k_1$	0.0067
⋮	⋮	⋮	⋮		⋮	⋮
$k_n$	0.0010	$k_n$	0.0011		$k_n$	0.0012

\*Probability values for illustration

# Problem statement

- How to choose the full key?

$k_0$	0.0065*	$k_0$	0.0071	.....	$k_0$	0.0070
$k_1$	0.0063	$k_1$	0.0068		$k_1$	0.0067
$\vdots$	$\vdots$	$\vdots$	$\vdots$		$\vdots$	$\vdots$
$k_n$	0.0010	$k_n$	0.0011		$k_n$	0.0012

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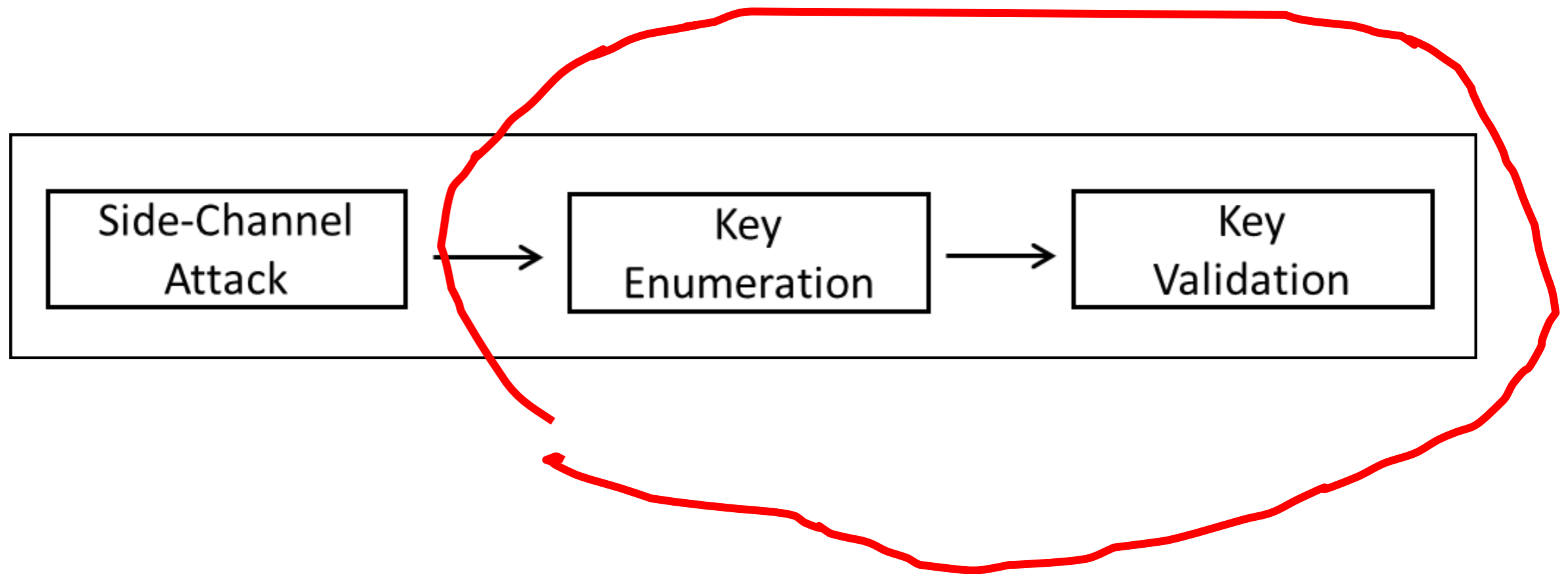
- Solution:
  - Key enumeration
  - Concurrent validation of full keys

\*Probability values for illustration



# Problem statement

Find the key after DPA as fast as possible. On a desktop.



DPA



lists of key chunk candidates with probabilities



full key?

attack



**enumeration**

in which **order** to  
brute force the full  
key candidates?

evaluation



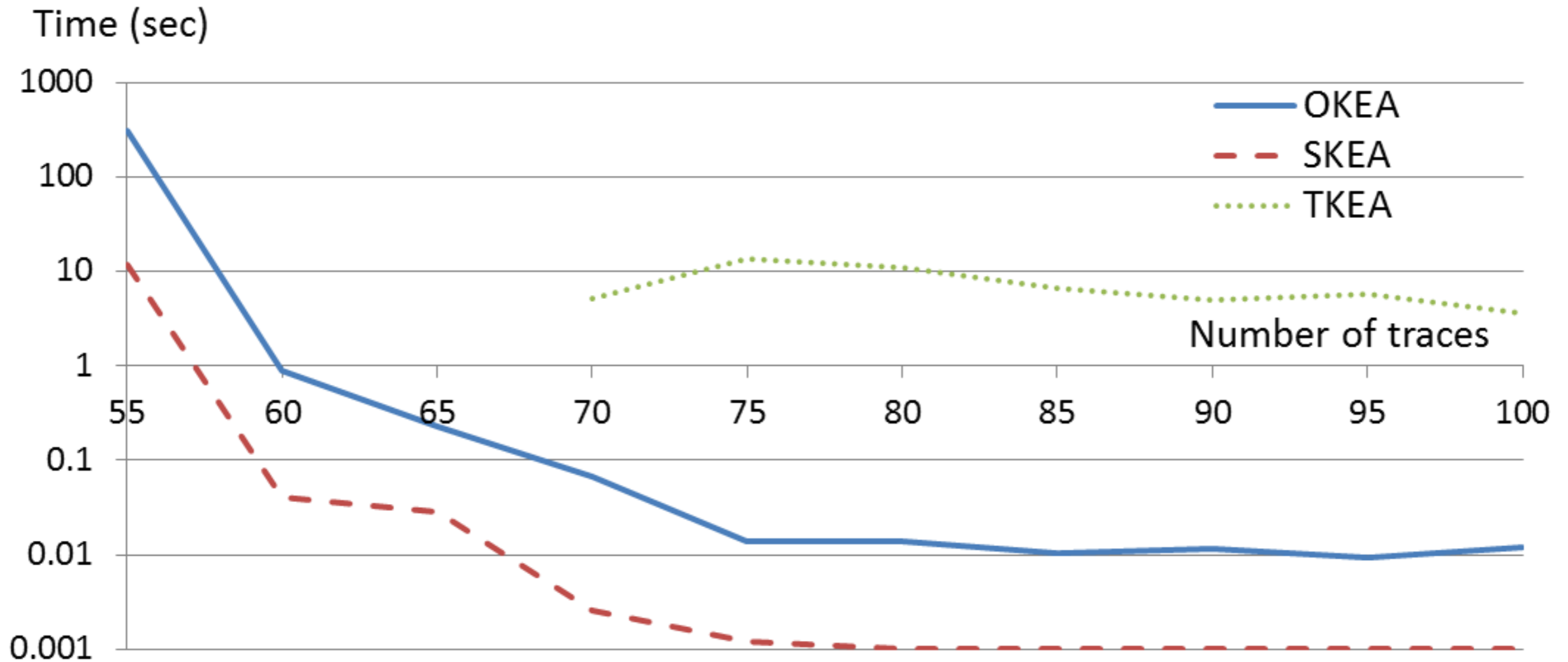
**rank estimation**

what is the **position**  
of the correct full key  
candidate?

# In this work so far...

- Benchmarked 3 enumeration algorithms:
  - Trivial Key Enumeration Algorithm (TKEA)
  - Optimal Key Enumeration Algorithm (OKEA) [SAC 2012]
  - **Score based Key Enumeration Algorithm (SKEA) [Marc Witteman] – new solution**
- Combined with key validation to get the full solution

# Time to find the key



# Trivial Key Enumeration Algorithm (TKEA)

$k_0$	0.0065*	$k_0$	0.0071	.....	$k_0$	0.0070
$k_1$	0.0063	$k_1$	0.0068		$k_1$	0.0067
$\vdots$	$\vdots$	$\vdots$	$\vdots$		$\vdots$	$\vdots$
$k_n$	0.0010	$k_n$	0.0011		$k_n$	0.0012

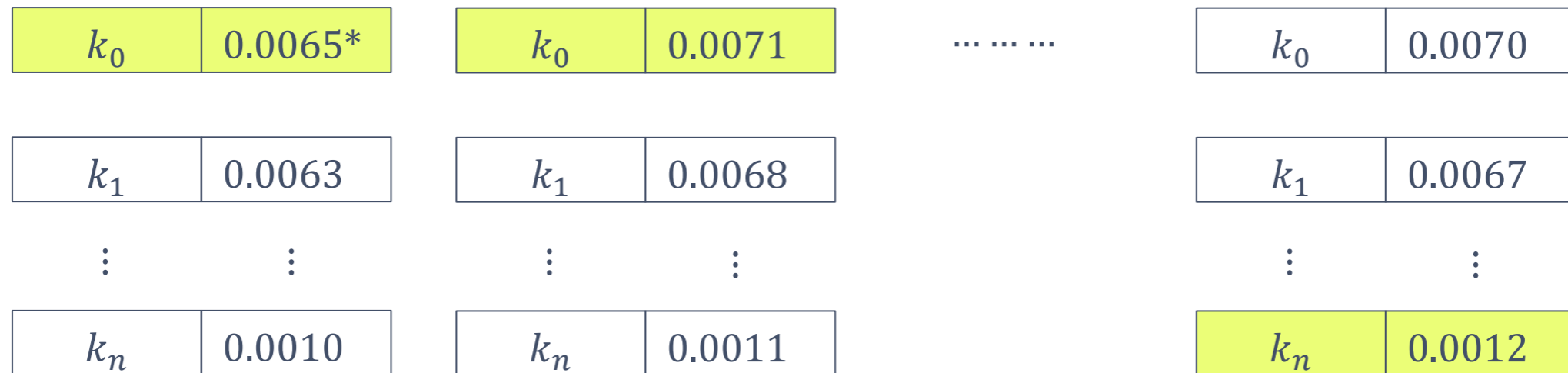
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# Trivial Key Enumeration Algorithm (TKEA)

$k_0$	0.0065*	$k_0$	0.0071	.....	$k_0$	0.0070
$k_1$	0.0063	$k_1$	0.0068		$k_1$	0.0067
$\vdots$	$\vdots$	$\vdots$	$\vdots$		$\vdots$	$\vdots$
$k_n$	0.0010	$k_n$	0.0011		$k_n$	0.0012

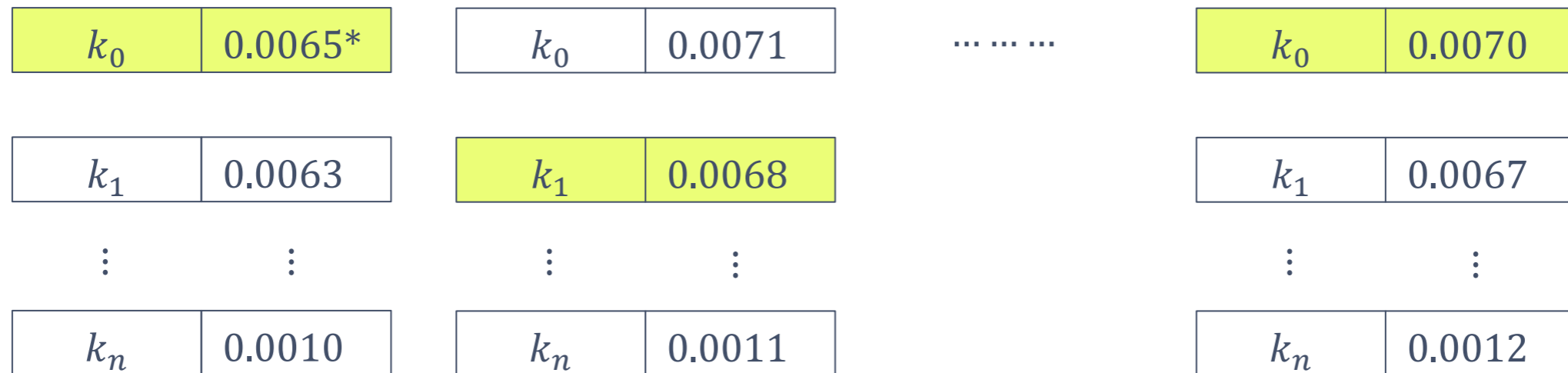
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# Trivial Key Enumeration Algorithm (TKEA)



\*Probability values for illustration

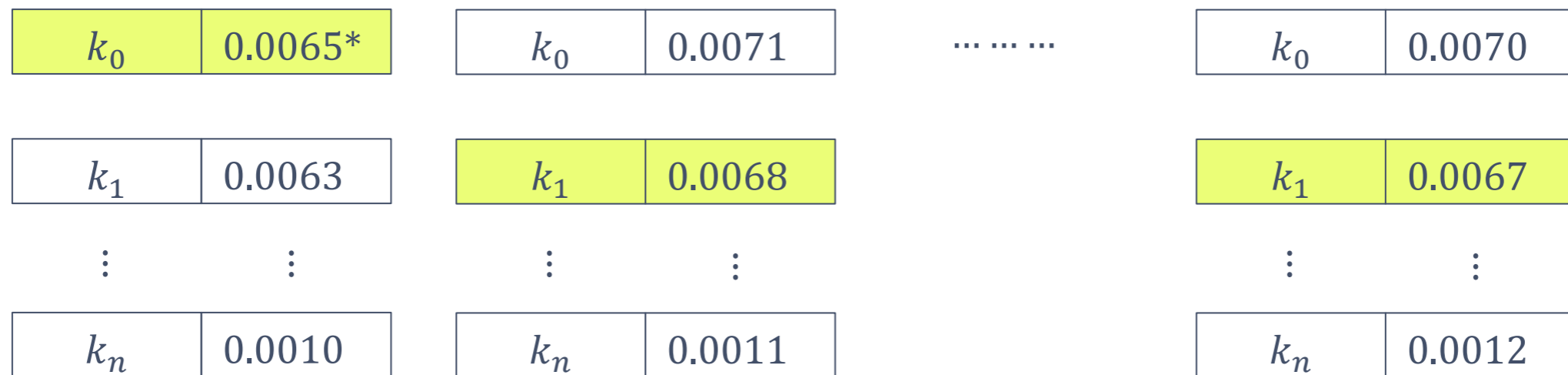
# Trivial Key Enumeration Algorithm (TKEA)



\*Probability values for illustration

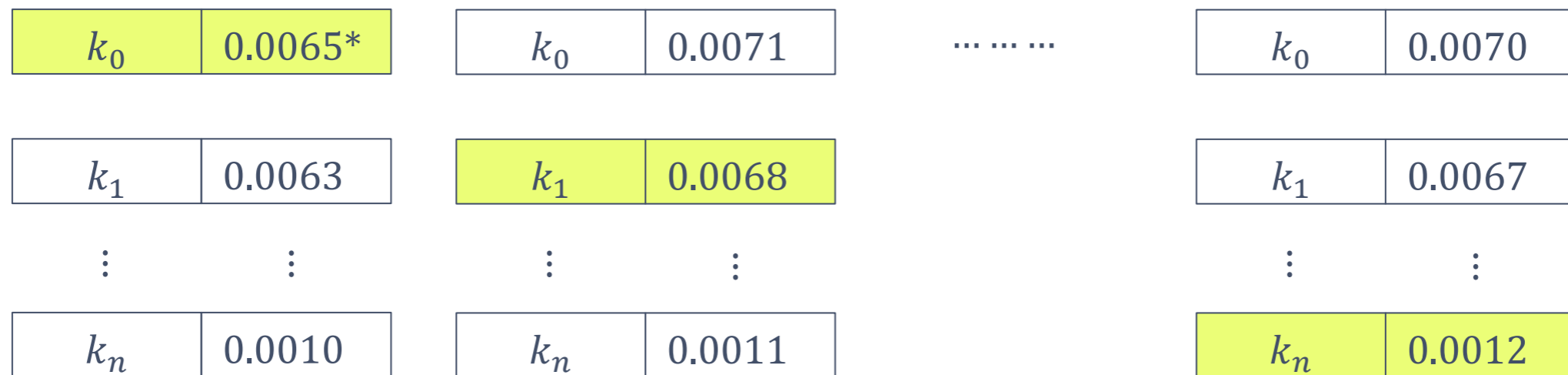


# Trivial Key Enumeration Algorithm (TKEA)




\*Probability values for illustration

# Trivial Key Enumeration Algorithm (TKEA)



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# Trivial Key Enumeration Algorithm (TKEA)

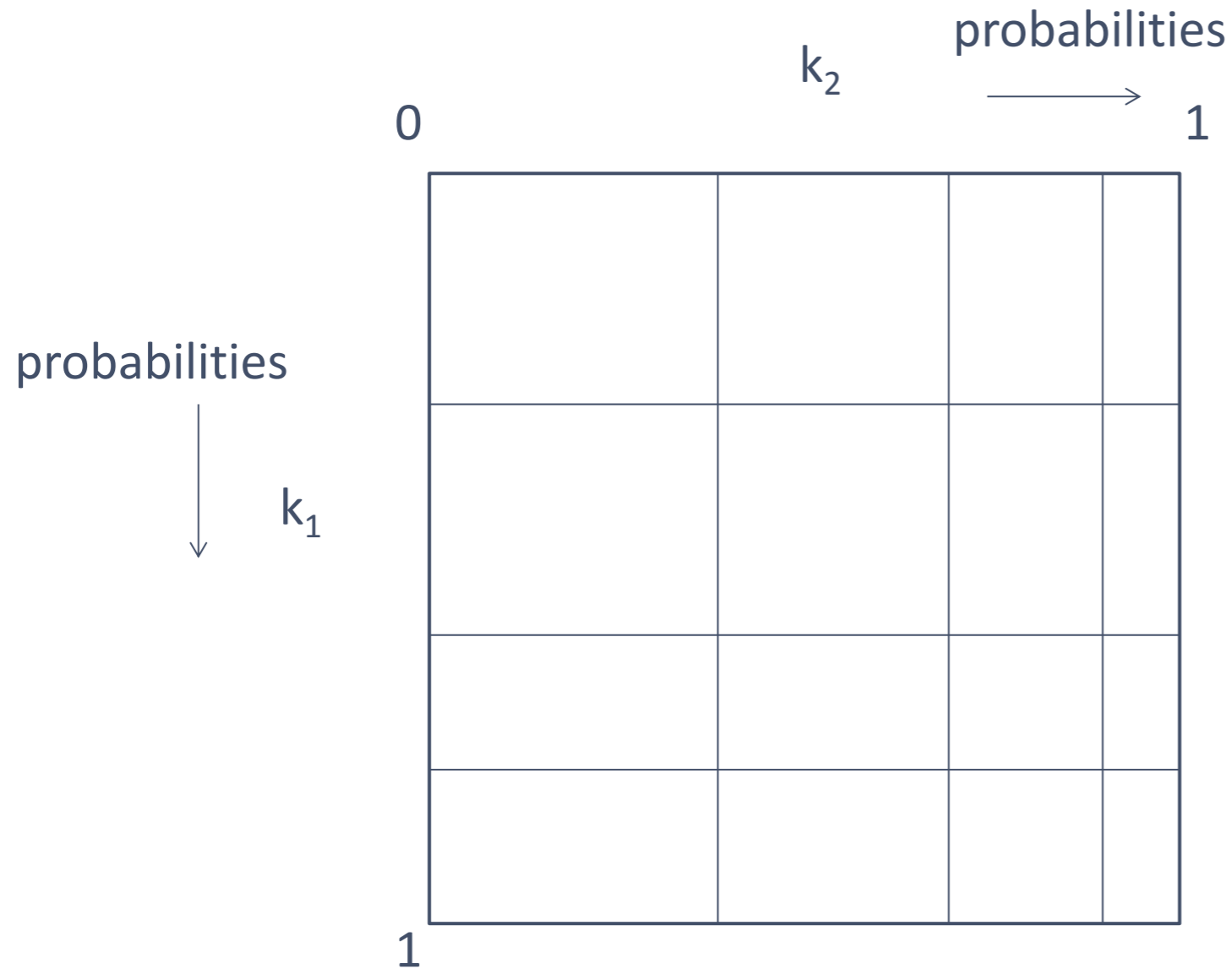
- Analysis:
  - Does not consume significant memory
  - Does not follow optimality at all
  - Has a very high throughput
  - Does not generate key within a feasible amount of time
- Above claims are based on our experimental results
- Not of our interest... 

# Optimal Key Enumeration Algorithm (OKEA) [SAC 2012]

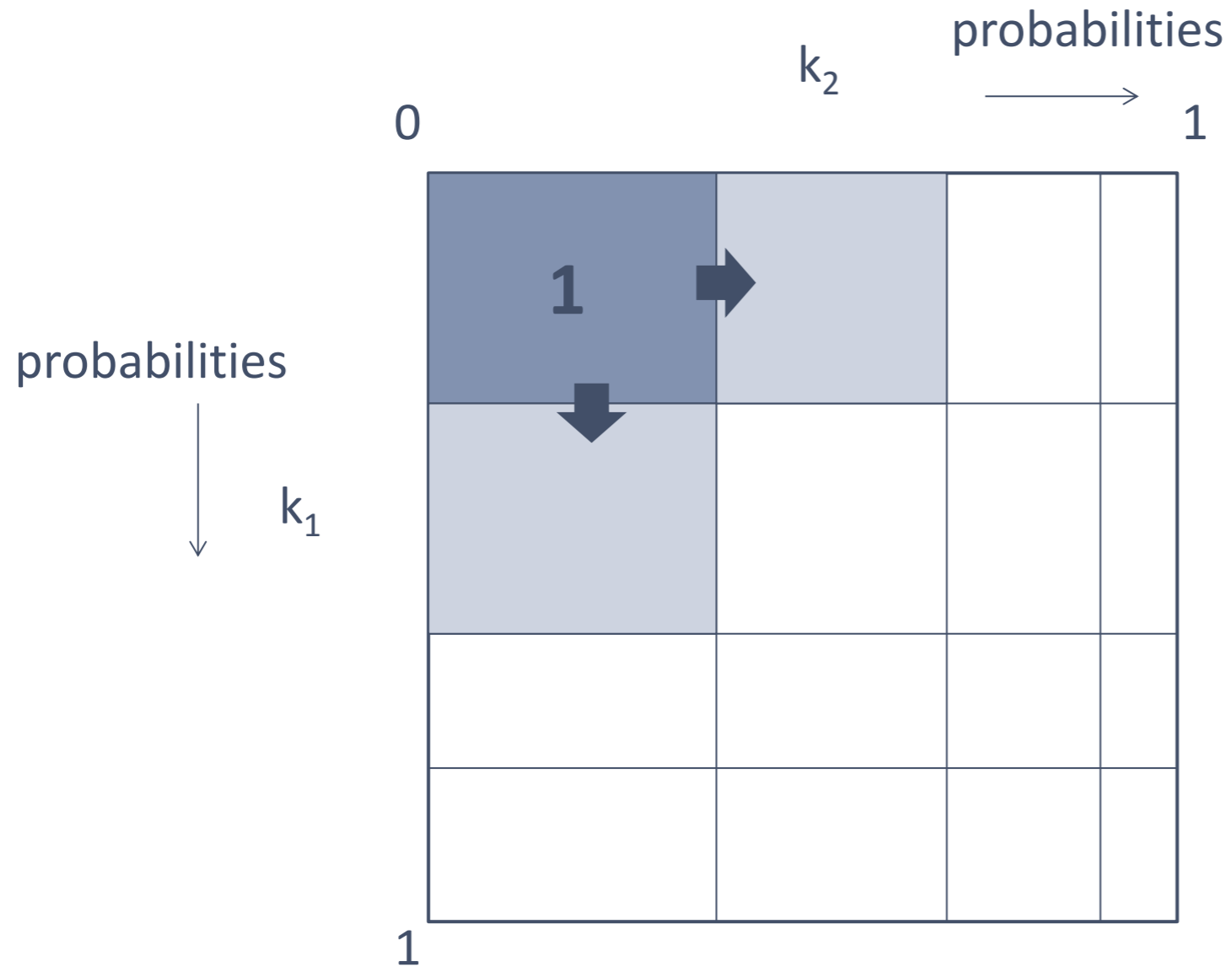


- Proposed by Nicolas Veyrat-Charvillon, Benoît Gérard, Mathieu Renaud and François-Xavier Standaert in 2012
- Requires probability based ranking of sub-key lists
- Let us discuss a simple case of merging two lists

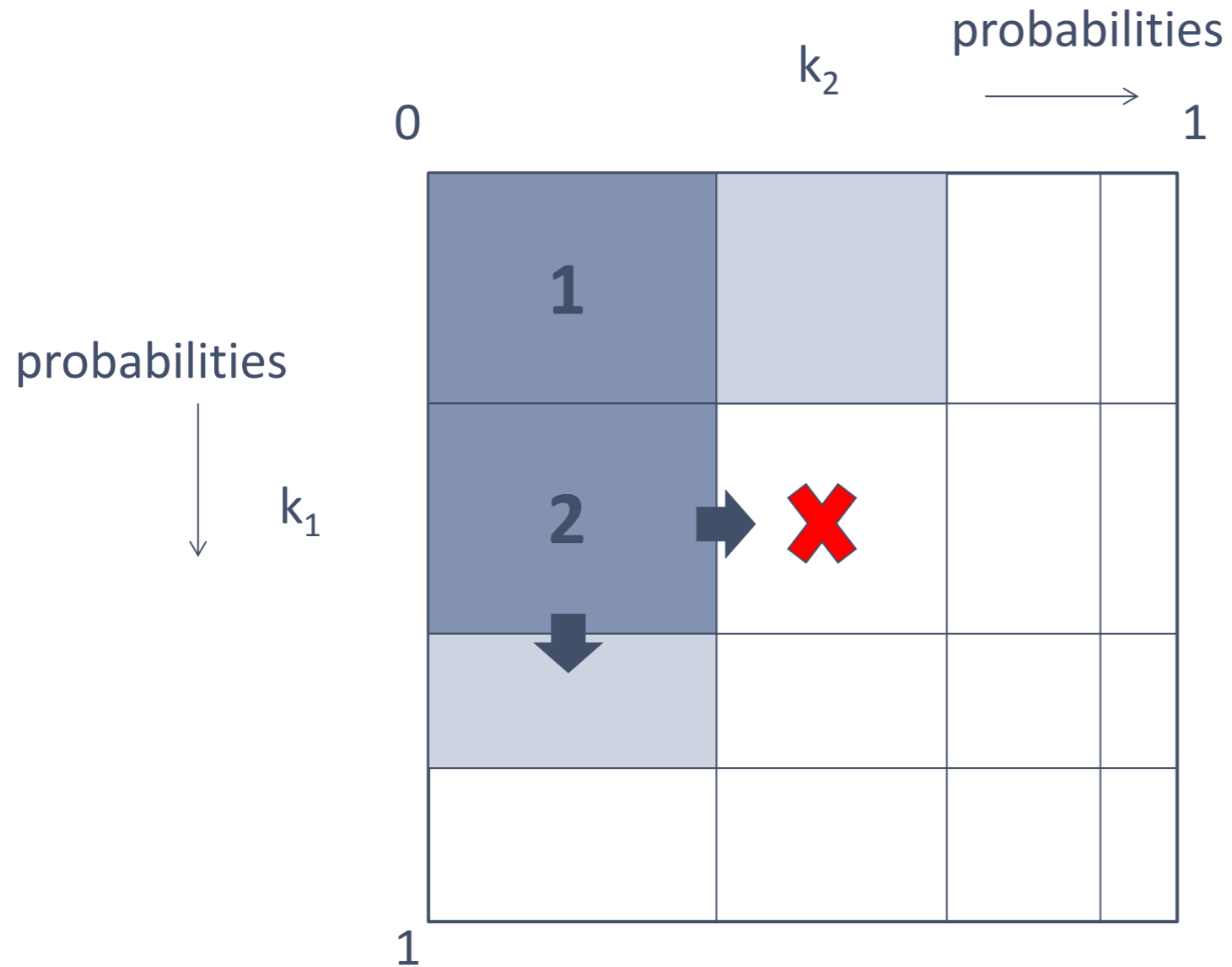
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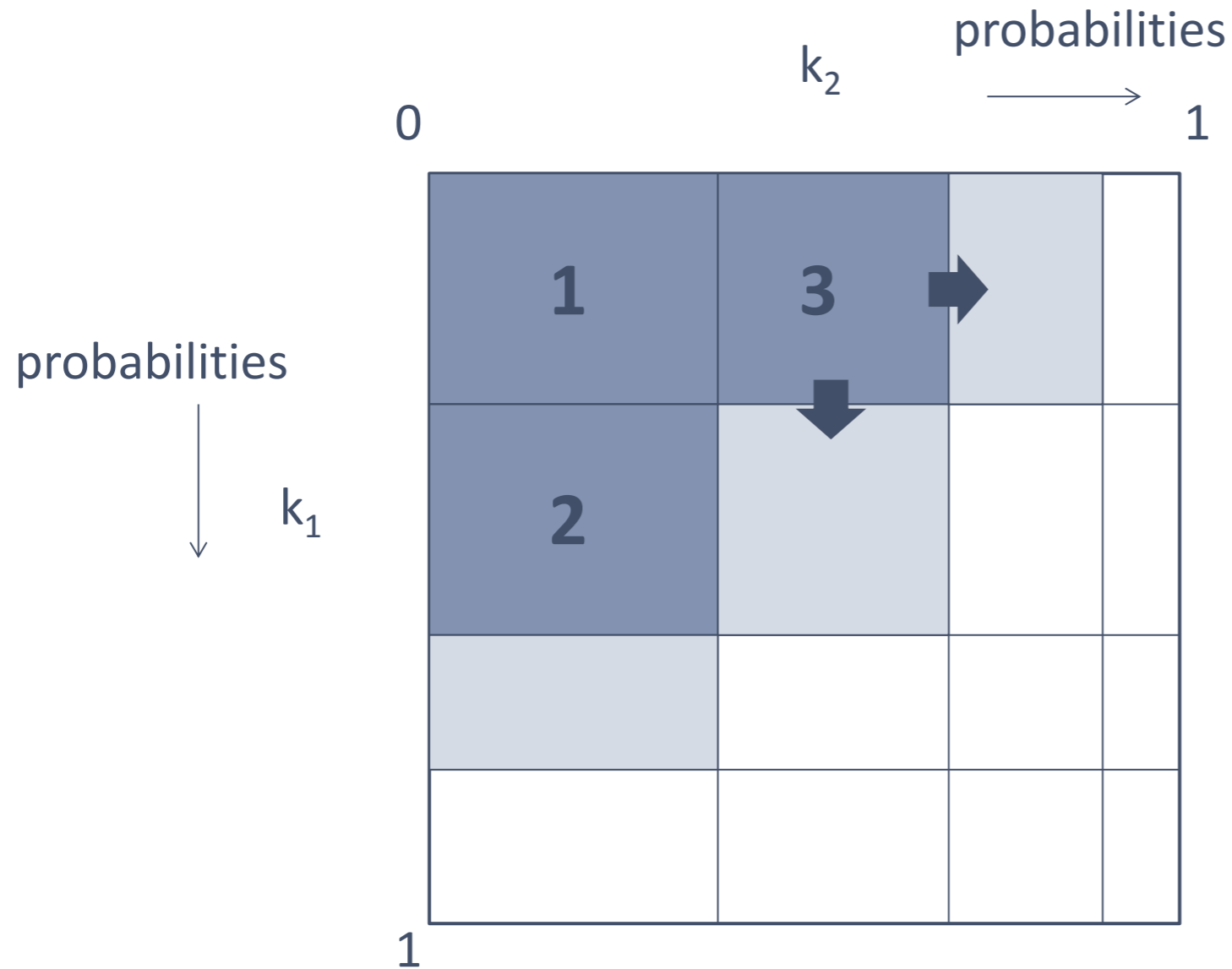
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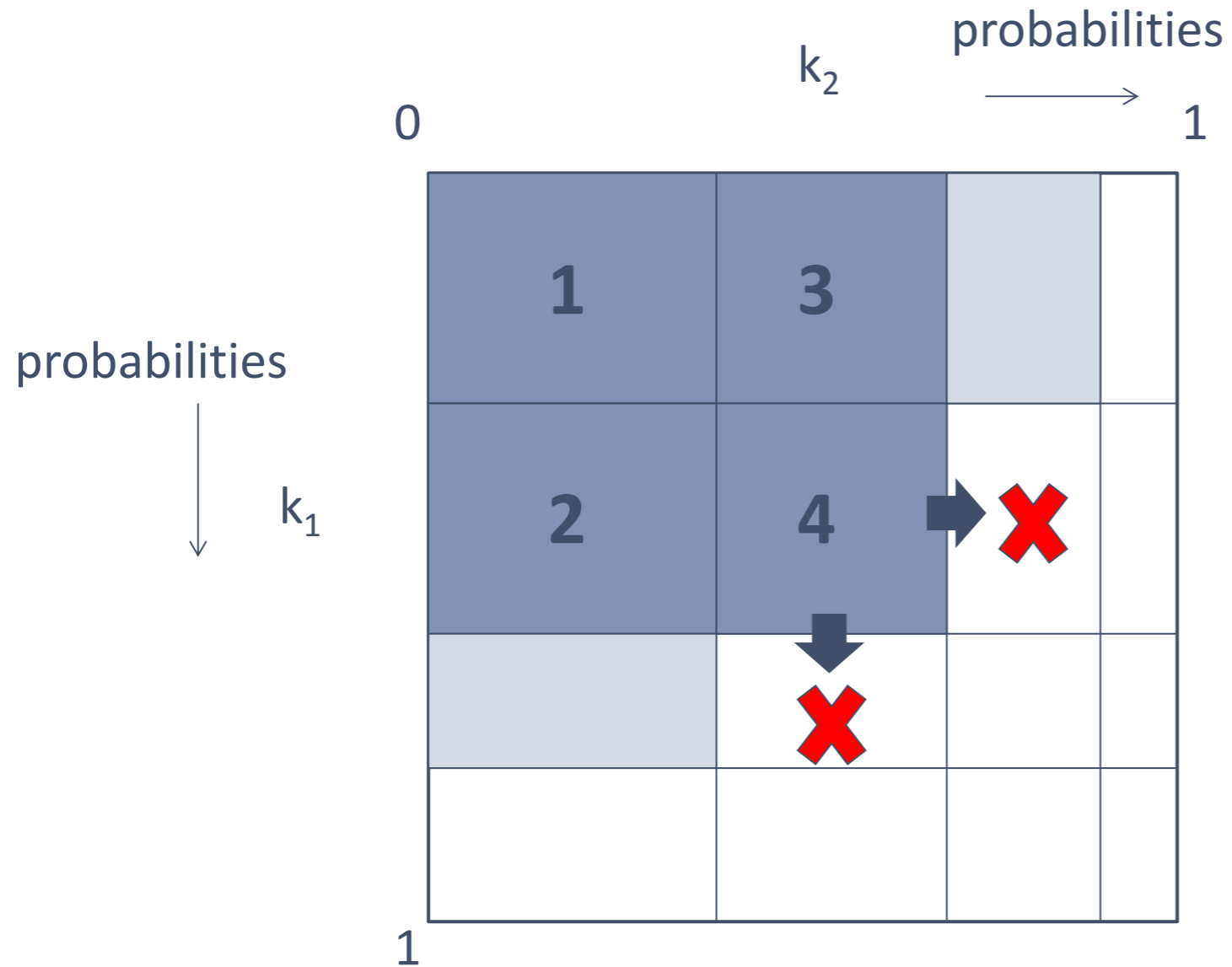


# Optimal Key Enumeration Algorithm (OKEA) [SAC 2012]

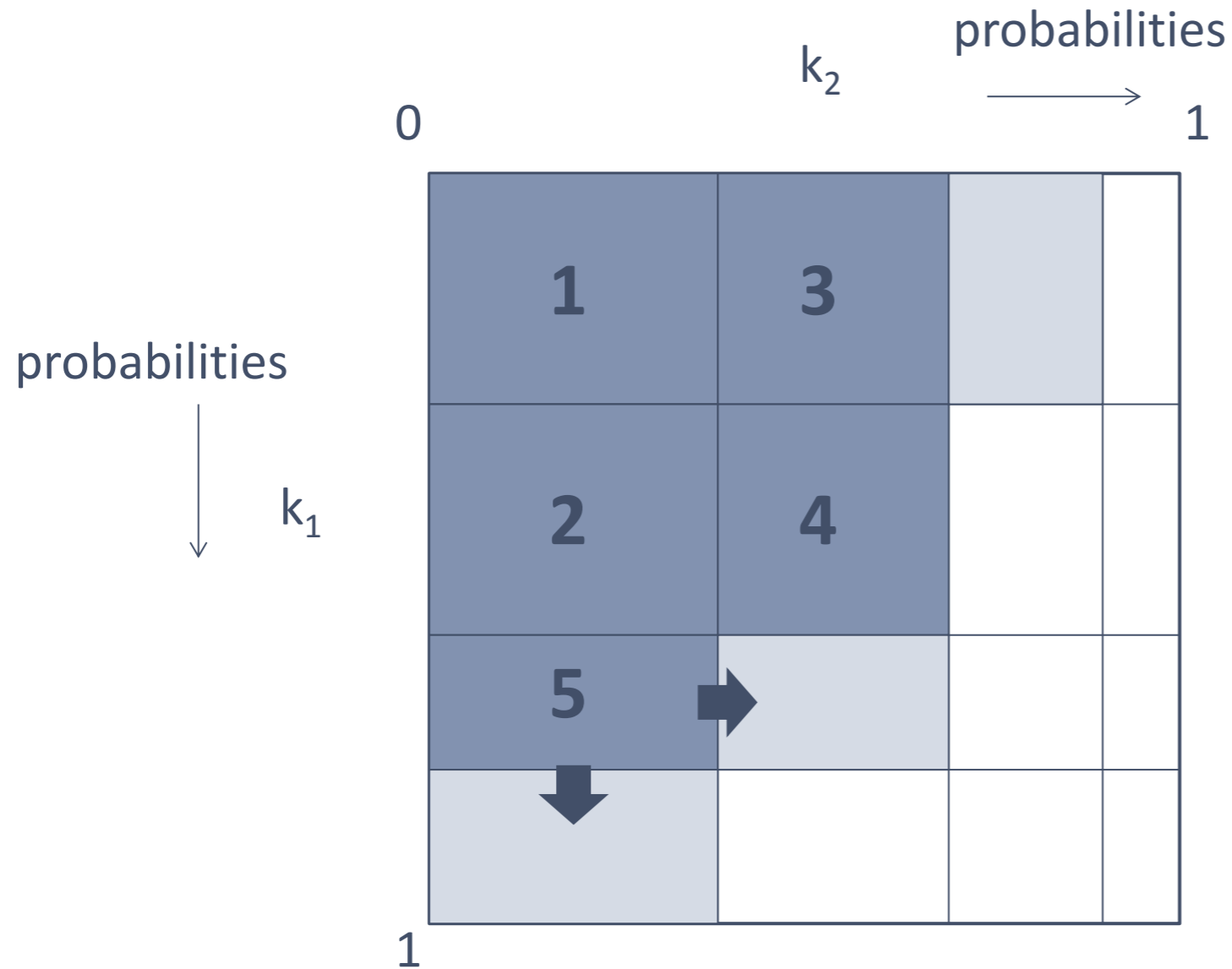




# Optimal Key Enumeration Algorithm (OKEA) [SAC 2012]

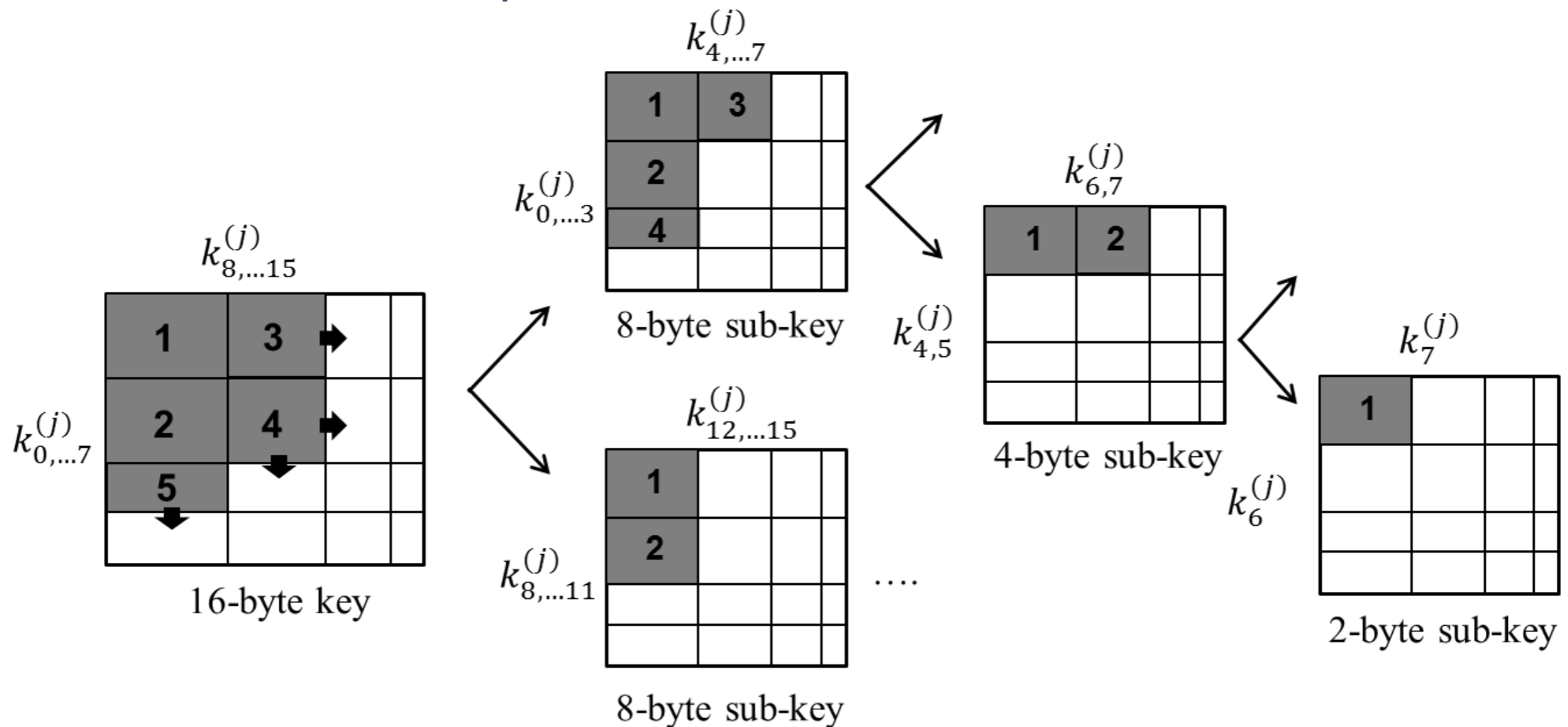


# Optimal Key Enumeration Algorithm (OKEA) [SAC 2012]



# Optimal Key Enumeration Algorithm (OKEA) [SAC 2012]

- Merging n lists:
  - merging two lists n-1 times
  - recursive decomposition



# OKEA Performance on CPU

## [SAC 2012]




Figures from the paper

<b># key candidates</b>	$2^{12}$	$2^{16}$	$2^{20}$	$2^{24}$	$2^{28}$	$2^{32}$	$2^{36}$	$2^{40}$
<b>Time</b>	0s	0.03s	0.55s	9.2s	163s	3130s	12h	221h
<b>Memory</b>	88KB	405KB	2.7MB	20MB	225MB	1.8GB	10GB	70GB
<b>Throughput keys/sec</b>	-	$2^{21}$	$2^{21}$	$2^{21}$	$2^{21}$	$2^{20}$	$2^{21}$	$2^{20}$

predictions!

A black arrow pointing from the text "predictions!" towards the rightmost two columns of the table above.

# Optimal Key Enumeration Algorithm (OKEA) [SAC 2012]

- Analysis:
  - Consumes large amount of memory
  - Strictly follows optimality
  - Has very low throughput
  - Takes large amount of time to generate the correct key
- Above claims are based on our experimental results
- Still not of our interest... 

# Score based Key Enumeration Algorithm (SKEA) [Marc Witteman]

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- Requires discrete score based ranking of sub-key lists (conversion e.g. by multiplication and rounding)
- Main contribution of this work
- Explained in the following slides at a toy example (4 key chunks)

# Score based Key Enumeration Algorithm (SKEA) [Marc Witteman]



Key part →

Score ↓

	0	1	2	3
0	5	5	4	5
1	2	3	2	4
2	1	1	0	3
3	0	1	0	2

Best cumulated score:  $5+5+4+5=19$

Next:  $5+5+4+4=18$

Enumeration is done by descending cumulated score

# Score based Key Enumeration Algorithm (SKEA) [Marc Wittenman]

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19

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

18

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

17

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

16

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2

5	5	4	5
2	3	2	4
1	1	0	3
0	1	0	2



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Key part →

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	0	1	2	3
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		Key part →			
		0	1	2	3
Score ↓	0	5	5	4	5
	1	2	3	2	4
	2	1	1	0	3
	3	0	1	0	2

Cumulated Maximum			
19	14	9	5

# Score based Key Enumeration Algorithm (SKEA) [Marc Witteman]



Key part →

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	0	1	2	3
0	5	5	4	5
1	2	3	2	4
2	1	1	0	3
3	0	1	0	2

Cumulated Maximum

19	14	9	5
----	----	---	---

Cumulated Minimum

3	3	2	2
---	---	---	---

# Score based Key Enumeration Algorithm (SKEA) [Marc Witteman]



Key part →

Score ↓

	0	1	2	3
0	5	5	4	5
1	2	3	2	4
2	1	1	0	3
3	0	1	0	2

Cumulated Maximum

19	14	9	5
----	----	---	---

Cumulated Minimum

3	3	2	2
---	---	---	---

Key part →

Cumulated score ↓

	0	1	2	3
19				
18				
17				
16				
15				
14				
13				
12				
11				
10				
9				
8				
7				
6				
5				
4				
3				
2				

Table shows for each key part for a given accumulated score

- Minimal index
- Maximal index

# Score based Key Enumeration Algorithm (SKEA) [Marc Witteman]



Key part →

Score ↓

	0	1	2	3
0	5	5	4	5
1	2	3	2	4
2	1	1	0	3
3	0	1	0	2

Cumulated Maximum

19	14	9	5
----	----	---	---

Cumulated Minimum

3	3	2	2
---	---	---	---

Key part →

Cumulated score ↓

	0	1	2	3
19	0 / 0			
18				
17				
16				
15				
14				
13				
12				
11				
10				
9				
8				
7				
6				
5				
4				
3				
2				

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# Score based Key Enumeration Algorithm (SKEA) [Marc Witteman]



Key part →

Score ↓

	0	1	2	3
0	5	5	4	5
1	2	3	2	4
2	1	1	0	3
3	0	1	0	2

Cumulated Maximum

19	14	9	5
----	----	---	---

Cumulated Minimum

3	3	2	2
---	---	---	---

Key part →

Cumulated score ↓

	0	1	2	3
19	0 / 0			
18	0 / 0			
17	0 / 0			
16	0 / 1			
15	0 / 2			
14	0 / 3	0 / 0		
13	0 / 3	0 / 0		
12	0 / 3	0 / 1		
11	0 / 3	0 / 1		
10	0 / 3	0 / 3		
9	0 / 3	0 / 3	0 / 0	
8	0 / 3	0 / 3	0 / 0	
7	1 / 3	0 / 3	0 / 1	
6	1 / 3	1 / 3	0 / 1	
5	2 / 3	2 / 3	1 / 3	0 / 0
4	3 / 3	2 / 3	1 / 3	1 / 1
3	3 / 3	2 / 3	2 / 3	2 / 2
2				3 / 3

Table shows for each key part for a given accumulated score

- Minimal index
- Maximal index

# Score based Key Enumeration Algorithm (SKEA) [Marc Witteman]

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- Analysis:
  - Consumes very limited constant amount of memory
  - Sub-optimal
  - Has very high throughput
  - Takes relatively less time to reach the correct key
- Above claims are based on our experimental results
- Satisfies most of the properties of our interest. 😊

# Experimental comparison

- Key enumeration algorithms under consideration work for all side-channel attacks that employ divide and conquer strategy.
- In order to evaluate their performance, we just need sorted sub-key lists.
- Executed a correlation coefficient based Differential Power Analysis (DPA) attack on ATmega implementation of AES-128.
- Yielded 16 sub-key lists sorted according to the correlation scores of the sub-key candidates.

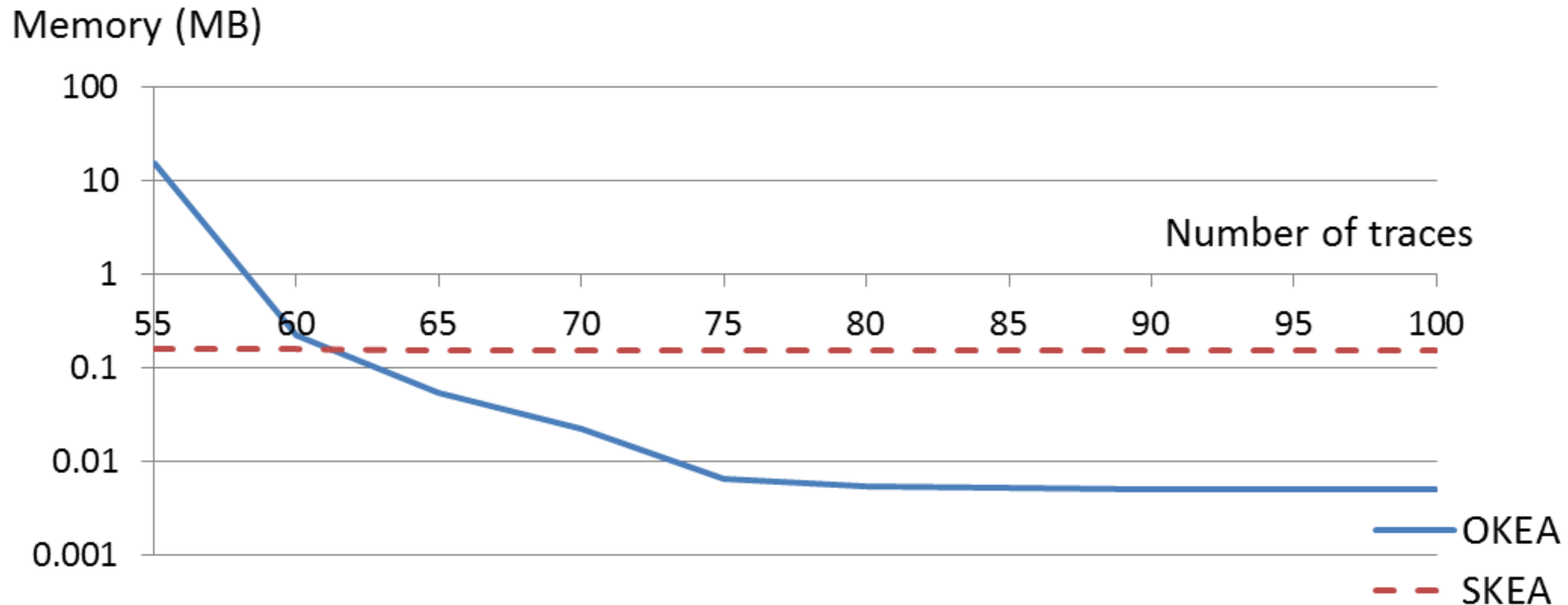


# Experimental Comparisons

- The attack was executed with varying the number of power traces 55, 60, ..., 100.
- Collected 12 cases for each number of power traces. 12 cases for 100 traces, 12 cases for 95 and so on.
- Executed the key enumeration algorithms on a common laptop (Intel core i5-2450M with 8 GB RAM and running 64-bit Windows 7).
- 1 hour was the time limit to find the correct key.
- Both OKEA and SKEA were able to find the correct key within 1 hour. TKEA could only find the key till 70 traces.

# Experimental Comparison

## Memory Consumption

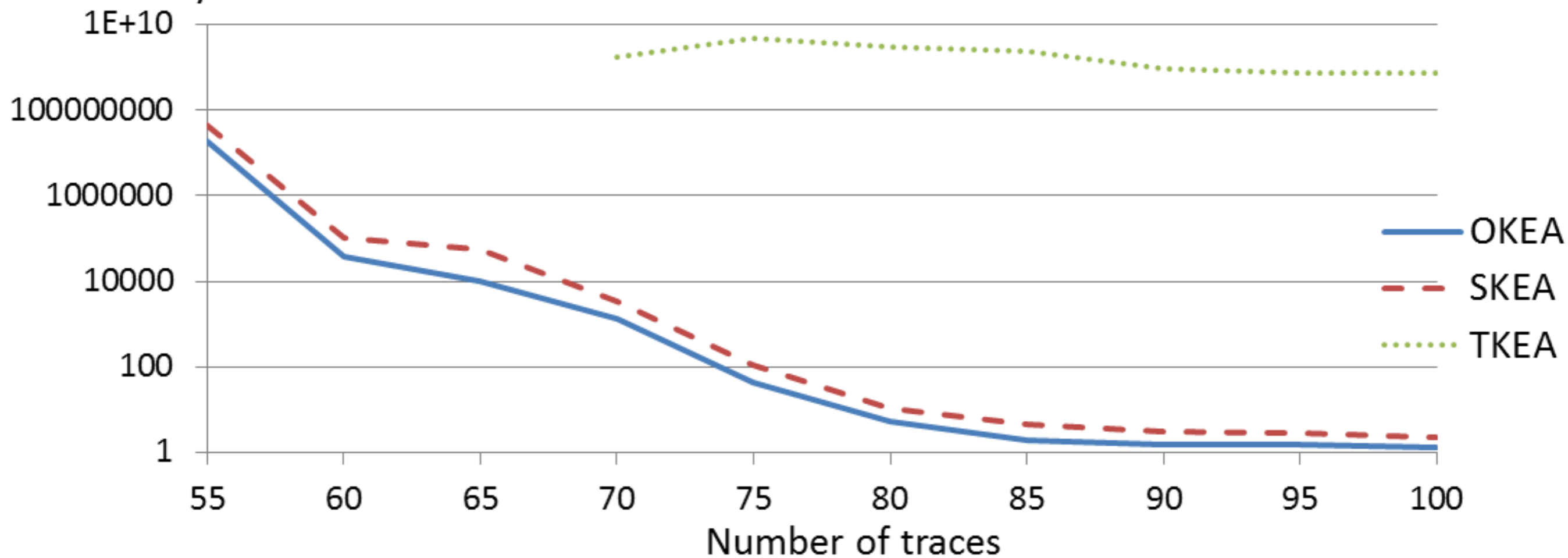


# Experimental Comparison



## Optimality

Rank of correct  
key

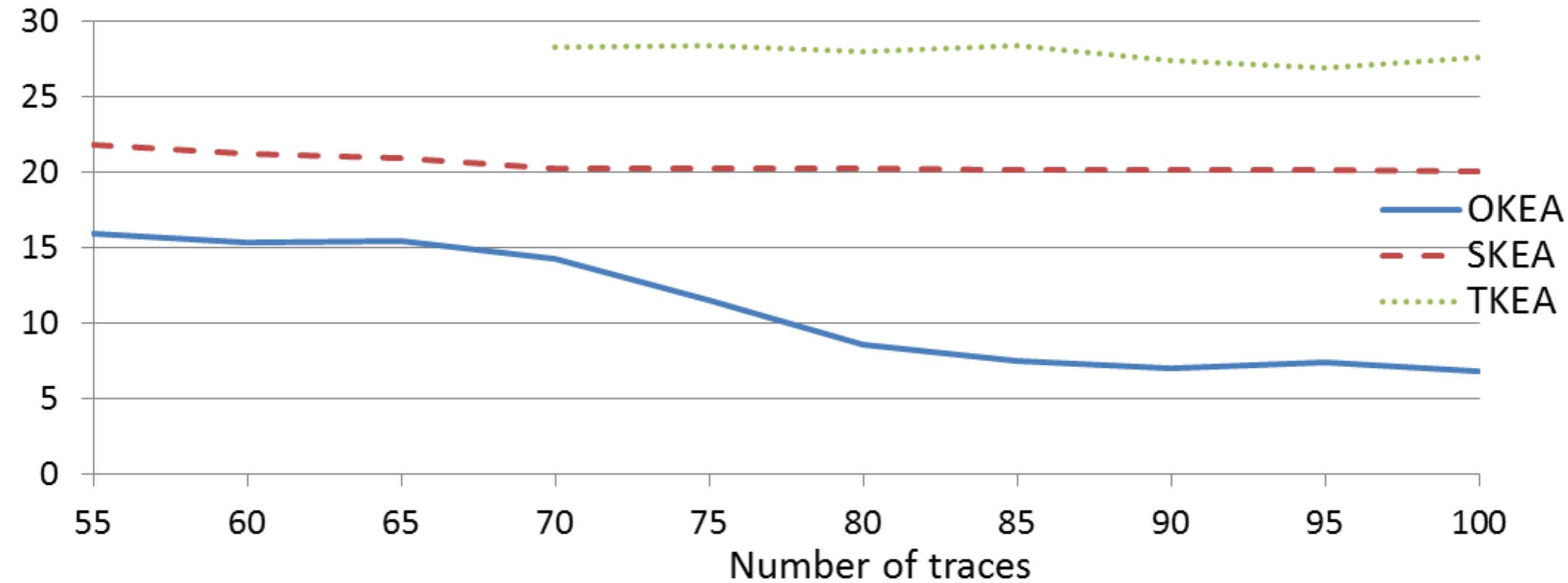


# Experimental Comparisons



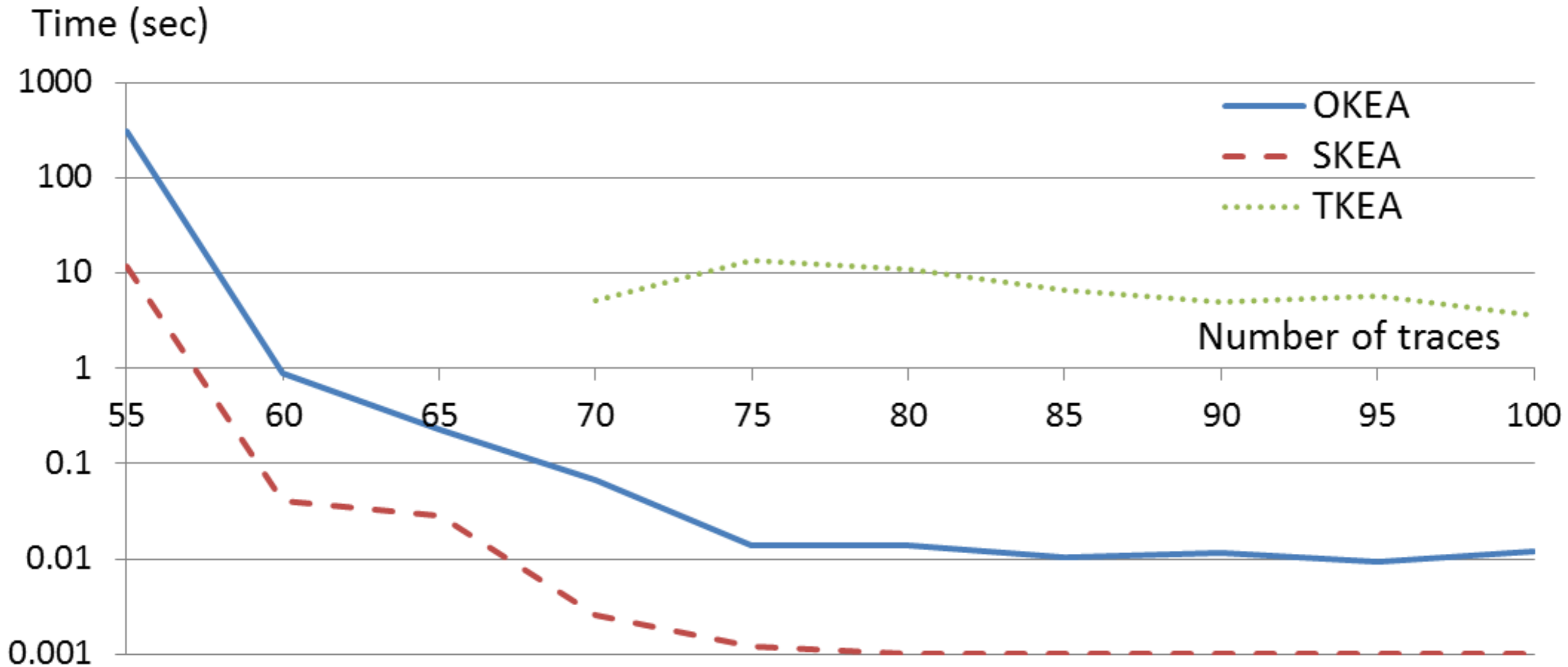
## Throughput

Throughput  
(keys/sec)  
in powers of 2



# Experimental Comparisons

## Time to find the Correct Key



# Discussion

- None of three key enumeration algorithms satisfy all the properties of our interest
- In contrast to both TKEA and OKEA, SKEA satisfies most of them
- **We propose to deploy SKEA as key enumeration algorithm**

# Key Validation

- Second step of key recovery mechanism
- Test whether the key generated by the key enumeration step is correct or not
- In order to validate a key, an attacker encrypts a particular plaintext with that key and matches the corresponding ciphertext with the one already computed with the stored secret key
- **Key validation rate should be compatible with key generation rate**

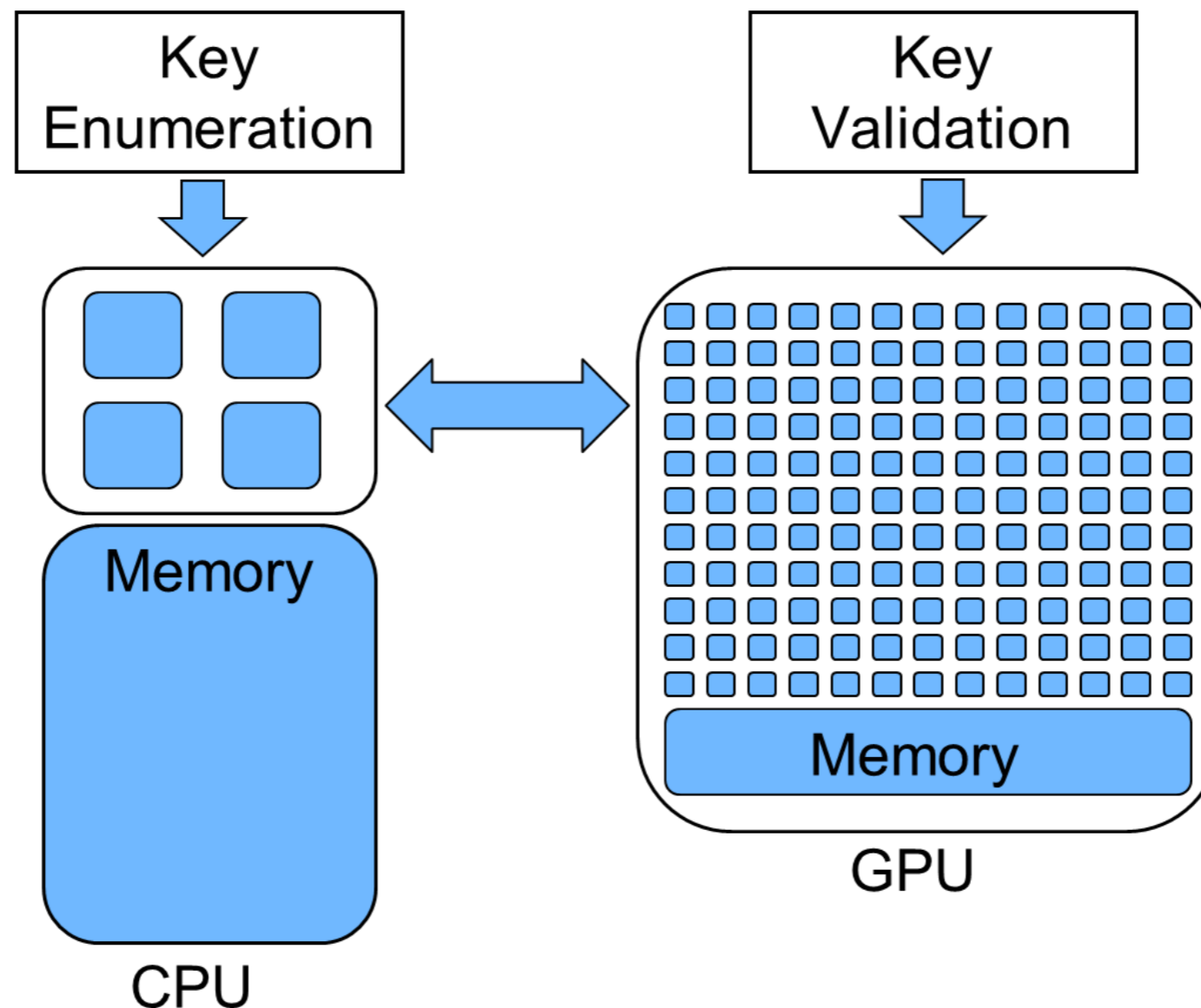
# Key Validation

- We propose to deploy SKEA. Maximum throughput of SKEA  $\cong 2^{23}$  *keys/sec*
- Throughput of key validation should at least be  $\cong 2^{23}$  *keys/sec*
- Utilize the immense parallel processing power of a GPU
- Implemented AES on an NVIDIA GPU using CUDA platform
- Achieved key validation rate of more than  $\cong 2^{23}$  *keys/sec*



# Proposed Solution

- Simultaneous execution of key enumeration (SKEA) on a CPU and key validation on a GPU.



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# Challenge your security

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